

# *Advanced Logic 2014–15*

## Assignment 1

(deadline: Friday, February 13)

1. Prove or disprove the universal validity of the following formulas:

- (a)  $\Box\bot \vee \Diamond\top$
- (b)  $\Diamond(p \vee q) \rightarrow \Diamond p \vee \Diamond q$
- (c)  $\Diamond p \wedge \Diamond q \rightarrow \Diamond(p \wedge q)$
- (d)  $\Box(p \rightarrow q) \rightarrow (\Diamond p \rightarrow \Box q)$
- (e)  $\Box(p \rightarrow q) \rightarrow (\Box p \rightarrow \Box q)$

2. If the formula  $\Box\bot$  is valid in a frame  $\mathcal{F} = (W, R)$ , then what do you know of its transition relation  $R$ ? Prove your answer.

Same question for  $\Diamond\top$ .

3. Let  $\mathbf{K}$  be the class of frames defined by:

$$\mathbf{K} = \{ (W, R) \mid \forall xyz \in W (Rxy \wedge Rxz \rightarrow Ryz \vee Rzy) \}$$

- (a) Show that  $\mathbf{K}$  is characterized by the formula  $\kappa$ :

$$\kappa = \Box(\Box p \rightarrow q) \vee \Box(\Box q \rightarrow p)$$

that is, prove that for all frames  $\mathcal{F}$ :

$$\mathcal{F} \in \mathbf{K} \quad \text{if and only if} \quad \mathcal{F} \models \kappa$$

- (b) Why does it follow from the previous item that also the formula

$$\Box(p \rightarrow \Diamond q) \vee \Box(q \rightarrow \Diamond p)$$

characterizes the class  $\mathbf{K}$ ?

(see page 2)

4. Consider the frames  $\mathcal{A} = (A, R_A)$ ,  $\mathcal{B} = (B, R_B)$ ,  $\mathcal{C} = (C, R_C)$ :

$$\begin{array}{ll} A = \{a, b, c, d\} & R_A = \{(a, b), (a, c), (b, a), (b, d)\} \\ B = \{1, 2, 3, 4\} & R_B = \{(1, 2), (1, 3), (1, 4), (2, 1)\} \\ C = \{\circ, \bullet\} & R_C = \{(\circ, \circ), (\circ, \bullet)\} \end{array}$$

(a) Draw the transition graphs for  $\mathcal{A}$ ,  $\mathcal{B}$  en  $\mathcal{C}$ .

Let further  $\mathcal{M}_{\mathcal{A}}$ ,  $\mathcal{M}_{\mathcal{B}}$ ,  $\mathcal{M}_{\mathcal{C}}$  be models based on  $\mathcal{A}$ ,  $\mathcal{B}$ , and  $\mathcal{C}$ , respectively, such that all propositional variables are false in all states.

- (b) Find a modal formula that *distinguishes* state  $a$  of model  $\mathcal{M}_{\mathcal{A}}$  from state 1 of model  $\mathcal{M}_{\mathcal{B}}$ ; that is, a formula  $\varphi$  such that  $\mathcal{M}_{\mathcal{A}}, a \models \varphi$  while  $\mathcal{M}_{\mathcal{B}}, 1 \not\models \varphi$ .
- (c) Why can we conclude from the previous item that there cannot exist a bisimulation  $Z \subseteq A \times B$  such that  $(a, 1) \in Z$ ?
- (d) Show that  $\mathcal{M}_{\mathcal{A}}, a \not\leftrightarrow \mathcal{M}_{\mathcal{C}}, \circ$ .
- (e) Can you find a modal formula which distinguishes state  $a$  in  $\mathcal{M}_{\mathcal{A}}$  from state  $\circ$  in  $\mathcal{M}_{\mathcal{C}}$ ?